

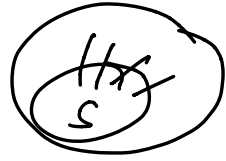
Announcements

- start recording
- scribe: work w/other scribe to produce one set of notes
- My OH: Mon 12:30-1:30, Thu 9-10
- TF hybrid section/OH: Mon 2-3, Wed 2-3, Wed 5-6, Thu 5:30-6:30
- PS 2 posted, including final project ideas problem
- If Zoom goes down, check Piazza
- sync whiteboard
- gather after class: link in chat
- jamboard link in chat

Agenda

- Recap/finish: proof of Cheeger
- Higher-order Cheeger

Recap: $\phi(S) \stackrel{\text{def}}{=} \frac{w(S, S^c)}{d(S)} = \frac{\sum_{a \in S, b \notin S} w_{ab}}{\sum_{a \in S} d(a)}$



$\phi(G) \stackrel{\text{def}}{=} \min_{S: d(S) \leq d(V)/2} \phi(S)$ Conductance of G

Thm: $\frac{\nu_2}{2} \leq \phi(G) \leq \sqrt{2\nu_2}$
 ↑
 Cheeger's Inequality

in a connected graph $\phi(G) > \frac{1}{n\Delta}$
 unweighted

Proof of Cheeger's Inequality

Goal: Given $\vec{y} \perp \vec{d}$ s.t. $\frac{\vec{y}^T L \vec{y}}{\vec{y}^T D \vec{y}} \leq \rho$ (e.s. $\vec{y} = D^{1/2} \vec{v}_2$, $\rho = \nu_2 \leq 2\phi(G)$)

"round" \vec{y} to obtain a set S (ie. a $\{0,1\}$ vector $\mathbb{1}_S$)

s.t. $\frac{w(S, S^c)}{\min\{d(S), d(V-S)\}} \leq \sqrt{2\rho} \leq 2\sqrt{2\phi(G)}$
 ↑
 if $\rho = \nu_2$

Sort y $y^{(1)} \leq \dots \leq y^{(n)}$



will take $S = \{a: y^{(a)} \leq t\}$ for some t

In fact, will define a distribution on thresholds t

$$\text{st. } \underline{E_{\tau} [\omega(\partial S_{\tau})]} \leq \underline{\sqrt{2\rho} \cdot E_{\tau} [\min \{d(S_{\tau}), d(V-S_{\tau})\}]}$$

Step 1: center \vec{y}

- $\underline{\vec{z}} = \underline{\vec{y}} - s\vec{1}$ for appropriate s

so that $\sum_{a: z(a) < 0} d(a) \leq \frac{d(V)}{2}$ and $\sum_{a: z(a) > 0} d(a) \leq \frac{d(V)}{2}$

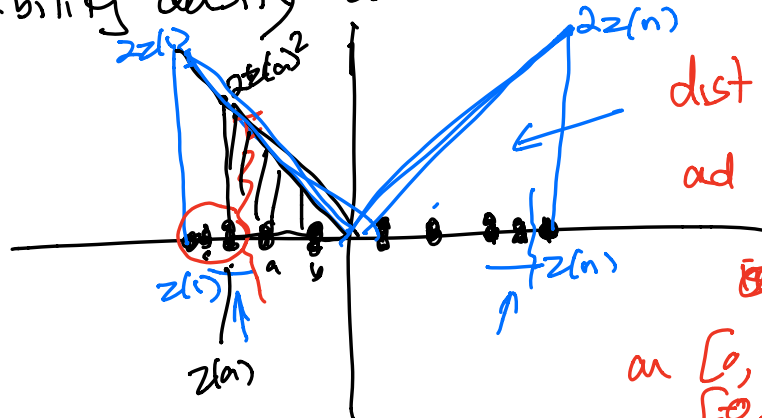
$$\frac{\vec{z}^T L \vec{z}}{\vec{z}^T D \vec{z}} \leq \rho$$

- Assume wlog $z(1)^2 + z(n)^2 = 1$

- $S_{\tau} \stackrel{\text{def}}{=} \{a : z(a) \leq \tau\}$

Step 2: distribution on τ

Probability density at $t = 2|t|$



dist of $t^2 | t > 0$
and dist of $t^2 | t < 0$

are uniform
on $[0, z(n)^2]$ and
 $[0, z(1)^2]$

Claim: $E_z \left[\min \{ d(S_\tau), d(V-S_\tau) \} \right] = \underline{z^T D z}$

Pf: LHS = $\sum_a \underline{d(a)} \cdot \Pr [a \text{ is in 'smaller' of } S_\tau \text{ and } V-S_\tau]$

centering
 $= \sum_a d(a) \cdot \Pr [z \text{ is between } z(a) \text{ and } 0]$
 $= \sum_a d(a) \cdot z(a)^2$
 $= z^T D z$

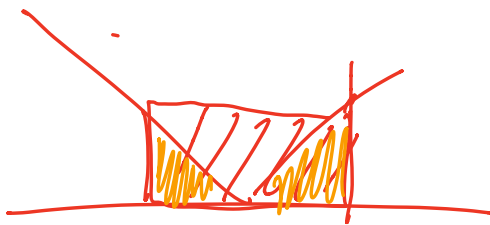
Claim: $E_z [w(\partial S_\tau)] \leq \sqrt{2} \sqrt{z^T L z} \cdot \sqrt{z^T D z}$

Pf: LHS = $\sum_{(a,b) \in E} w_{ab} \cdot \Pr [z \text{ is between } z(a) \text{ and } z(b)]$

$$\leq \sum_{(a,b) \in E} w_{ab} \cdot |z(b) - z(a)| \cdot (|z(b)| + |z(a)|)$$

$$\leq \sqrt{\sum_{(a,b) \in E} w_{ab} \cdot (z(b) - z(a))^2} \cdot \sqrt{\sum_{(a,b) \in E} w_{ab} \cdot (|z(b)| + |z(a)|)^2}$$

$$\leq \sqrt{z^T L z} \cdot \sqrt{2 z^T D z}$$



$$u_{ab} = \sqrt{w_{ab}} \cdot |z(b) - z(c)|$$

$$v_{ab} = \sqrt{w_{ab}} \cdot (|z(b)| + |z(c)|)$$

Cauchy-Schwarz Inequality

v1: for every two vectors $\vec{u}, \vec{v} \in \mathbb{R}^n$

$$|\langle \vec{u}, \vec{v} \rangle| \leq \|\vec{u}\| \cdot \|\vec{v}\|$$

v2: for all $u_1, \dots, u_n, v_1, \dots, v_n \in \mathbb{R}$

$$\left(\sum_{i=1}^n u_i v_i \right)^2 \leq \left(\sum_{i=1}^n u_i^2 \right) \cdot \left(\sum_{i=1}^n v_i^2 \right)$$

G undirected

Higher-Order Cheeger

Fact: $\nu_k = 0 \iff G$ has at least k connected components

Pf: 1) If G has exactly k connected components

$$N = \begin{pmatrix} N_1 & 0 & 0 & 0 \\ 0 & N_2 & 0 & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & N_k \end{pmatrix}$$

$$\frac{x^T N x}{x^T x} = \frac{y^T L y}{y^T D y}$$

$$y = D^{1/2} x$$

$N_i =$ normalized Laplacians for components
eigenvalues $(N) =$ union w/mult. of e-values of N_i

0 has mult. k

2) $\nu_k = 0 \iff \dim(\ker(N)) \geq k \iff \dim(\ker(L)) \geq k$
High-Order Cheeger: $\ker(L) = \text{span}(1_{S_1}, \dots, 1_{S_k})$

$$\nu_k \approx 0$$

\iff

G "close" to having $\geq k$ components
 $S_i =$ connected components

\exists disjoint $S_1, \dots, S_k \subseteq V$

s.t. $\phi(S_1), \dots, \phi(S_k)$ "small"

$$\phi_k(G) \stackrel{\text{def}}{=} \min_{S_1, \dots, S_k \text{ disjoint}} \max \{ \phi(S_1), \dots, \phi(S_k) \}$$

Thm: $\frac{\lambda_k}{2} \leq \phi_k(G) \leq \text{poly}(k) \cdot \sqrt{\lambda_k}$

Exercise for breakout s: try to reconstruct

proof that $\frac{\lambda_k}{2} \leq \phi_k(G)$

Let $\beta = \phi_k(G)$

Have disjoint sets S_1, \dots, S_k

s.t. $\phi(S_i) \leq \beta$ for each i

||

$$\frac{\mathbf{1}_{S_i}^T L \mathbf{1}_{S_i}}{\mathbf{1}_{S_i}^T D \mathbf{1}_{S_i}} \leq \beta$$

Q Let $S = \text{span}(\mathbf{1}_{S_1}, \dots, \mathbf{1}_{S_k})$

By C-F

$$\lambda_k \leq \max_{\substack{x \in S \\ x \neq 0}} \frac{x^T L x}{x^T D x} \stackrel{?}{\leq} 2\beta$$

↑
uses $\mathbf{1}_{S_1}, \dots, \mathbf{1}_{S_k}$
are disjointly supported

Proof idea for $\Phi_k(G) \leq \text{poly}(k) \cdot \sqrt{2k}$

Assume G d -regular and unweighted for simplicity

$$v \in \mathbb{R}^V$$

Let $F_1, \dots, F_k : V \rightarrow \mathbb{R}$ be

k orthonormal vectors s.t.

$$F_i^T N F_i = \frac{F_i^T L F_i}{d} \leq \rho \quad (\text{e.g. } \rho = 2k)$$

$\underbrace{\hspace{10em}}_{\rho_i}$

Define $F : V \rightarrow \mathbb{R}^k$ $F(a) = \begin{pmatrix} F_1(a) \\ \vdots \\ F_k(a) \end{pmatrix}$ cf. Hall's graph drawing

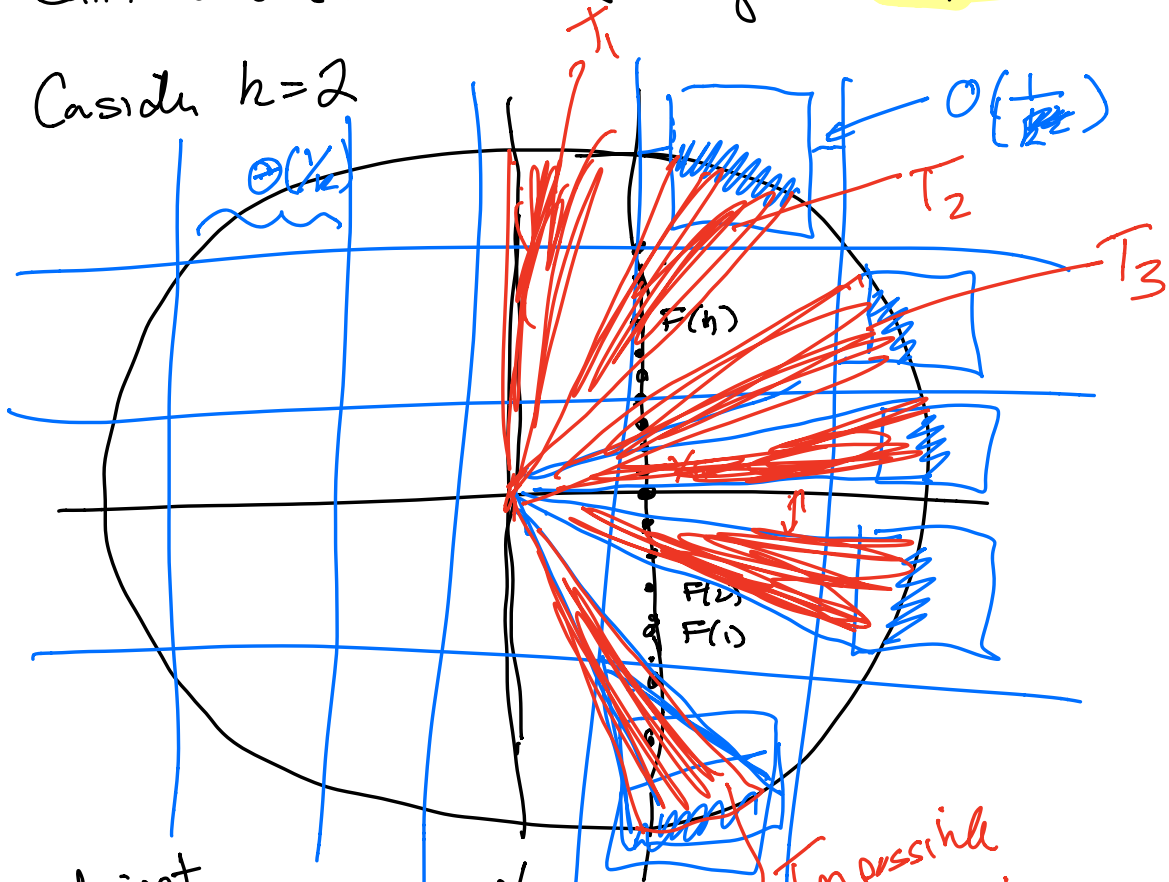
Know $\frac{\sum_{(a,b) \in E} \|F(a) - F(b)\|^2}{d} \leq \sum_{i=1}^k \rho_i = k \cdot \rho$

Also: $\sum_{a \in V} \|F(a)\|^2 = k$

How to "round" and get sets S_1, \dots, S_k ?

Can't round each F_i separately. (why?)

Consider $k=2$



\Rightarrow disjoint sets $T_1, \dots, T_m \subseteq V$ s.t.

T_m possible but $m \gg k$

- $\sum_{a \in T_i} \|F(a)\|^2 \geq k \cdot (1 - O(1/k))$

- $\forall i \sum_{a \in T_i} \|F(a)\|^2 \leq 1 + O(1/k)$

- $\forall i \neq j, a \in T_i, b \in T_j \left\| \frac{F(a)}{\|F(a)\|} - \frac{F(b)}{\|F(b)\|} \right\| \geq \frac{1}{poly(k)}$

$$\sqrt{2 - 2 \cos(\text{angle between } F(a) \text{ and } F(b))}$$

⇒ disjoint sets A_1, \dots, A_k s.t.

$$\bullet \forall i \sum_{a \in A_i} \|F(a)\|^2 \geq \frac{1}{2}$$

$$\bullet \forall i \neq j \ a \in A_i, b \in A_j \ \left\| \frac{F(a)}{\|F(a)\|} - \frac{F(b)}{\|F(b)\|} \right\| \geq \frac{1}{\text{poly}(k)}$$

$$y_i(a) = \begin{cases} \|F(a)\| & \text{if } a \in A_i \\ \text{interpolate to a.w.} & \\ 0 & \text{if } a \text{ is "far" from } A_i \end{cases}$$

"localization"

⇒ disjointly supported, nonnegative vectors y_1, \dots, y_k (on A_i)

$$\text{s.t. } \frac{y_i^T N y_i}{y_i^T y_i} \leq \text{poly}(k) \cdot \frac{F^T N F}{F^T F} = \text{poly}(k) \cdot \rho$$

⇒ disjoint sets S_1, \dots, S_k s.t.

Cheeger-like rounding

$$\phi(S_i) \leq \sqrt{2 \frac{y_i^T N y_i}{y_i^T y_i}} \leq \text{poly}(k) \cdot \sqrt{\rho}$$